

Functional **Meta**-Programming in the Construction of Parallel Programs

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Objectives

- efficient parallel target programs
 - predefined, parameterized parallel patterns
 - meta-programming to avoid overhead

Objectives

- efficient parallel target programs
- short development time
 - automatic program generation
 - use like a library

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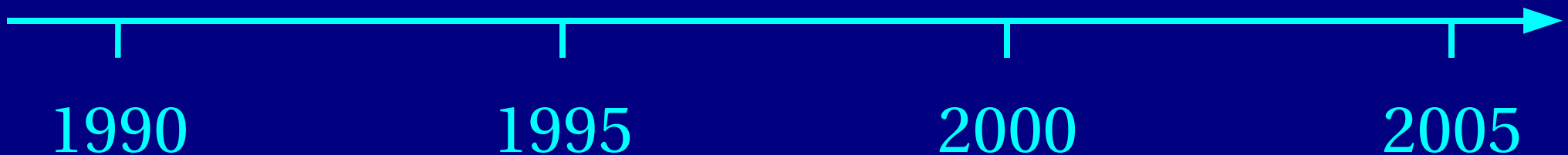
- efficient parallel target programs
- short development time
- appropriate for application programmer without experience of parallelism
 - focus on function
 - hide operational view

Objectives

- efficient parallel target programs
- short development time
- appropriate for application programmer without experience of parallelism
- software quality
 - semantic definition
 - cost model

Own and Related Approaches

- skeleton idea (Cole)
 - skeleton/ML compiler (Bratvold)
 - C macros
 - C++ skeletons (Kuchen)
 - SAT (Gorlatch)
 - HDC compiler
 - PSML comp. (Michaelson)
 - meta-progr.



New Approach (since 2002): Functional Meta-Programming (1)

Def. (meta-programming):

- analysis
- transformation
- generation

of object-programs by meta-programs

New Approach (since 2002): Functional Meta-Programming (2)

motivation for meta-programming:

- keep skeleton approach, but
 - (1) use existing compiler technology
 - (2) avoid administrative overhead

New Approach (since 2002): Functional Meta-Programming (3)

- CMPP 2002
 - domain-specific object-language
 - cost model
 - meta-language: Haskell

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- CMPP 2002
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 - meta-language: Haskell
- CMPP 2004
 - meta-language: MetaOCaml
 - + simple code-generation

Programming Layers

(1) domain-specific language for parallelism

- cost calculator
- SPMD code generator
- sequential/parallel task structure

Programming Layers

(1) domain-specific language for parallelism

(2) skeletons expressed in this language

- written by a parallelism expert
- handle the communications
- can exploit the parallel machine

Programming Layers

- (1) domain-specific language for parallelism
- (2) skeletons expressed in this language
- (3) application programming
 - based on skeletons
 - no knowledge in parallelism required

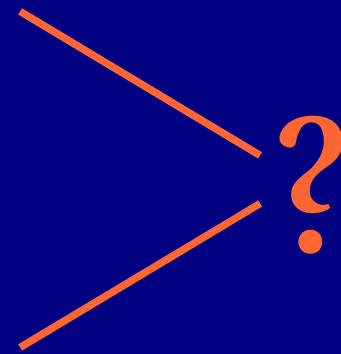
Design Alternatives for the parallel language

- small implementation effort

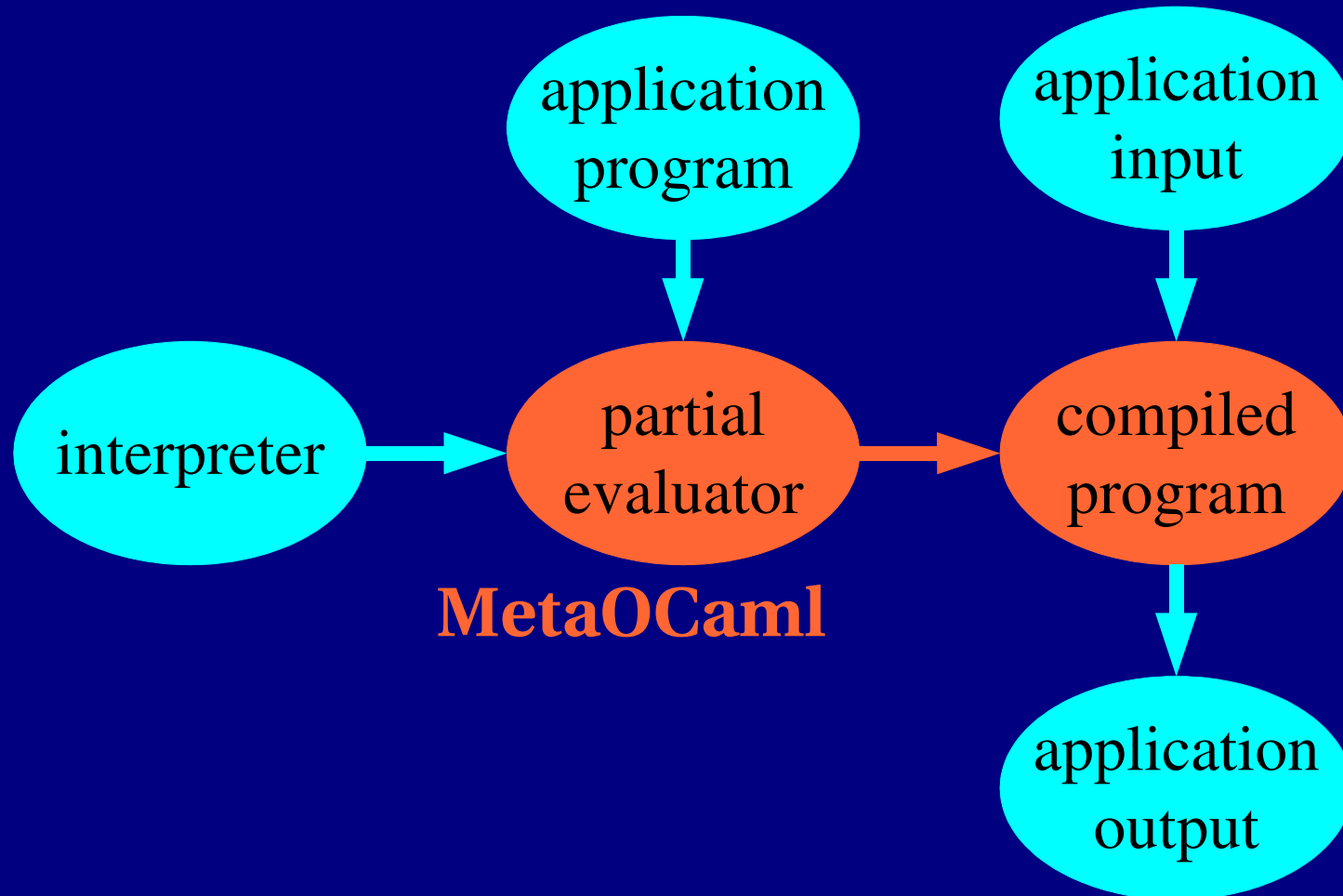
→ interpretation

- good efficiency

→ compilation



Interpretation + Partial Evaluation = Compilation



MetaOCaml

- developed by Walid Taha (Rice Univ.)
- based on Objective Caml
- run-time program specialization
- meta-programming extensions

Meta-Programming Extensions

brackets (`.< >.`): enclose object program part

```
# let a = .< 2*4 >. ;;  
val a : ('a, int) code = .<(2 * 4)>.
```

Meta-Programming Extensions

brackets (`.< >.`): enclose object program part

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# let a = .< 2*4 >. ;;  
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escape (`.~`): inserts object program part

```
# let b = .< 9 + .~a >. ;;  
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```

run (`.!`): executes object program

```
# let c = .!b ;;  
val c : int = 17
```

Use of Meta-Programming

one
meta-program

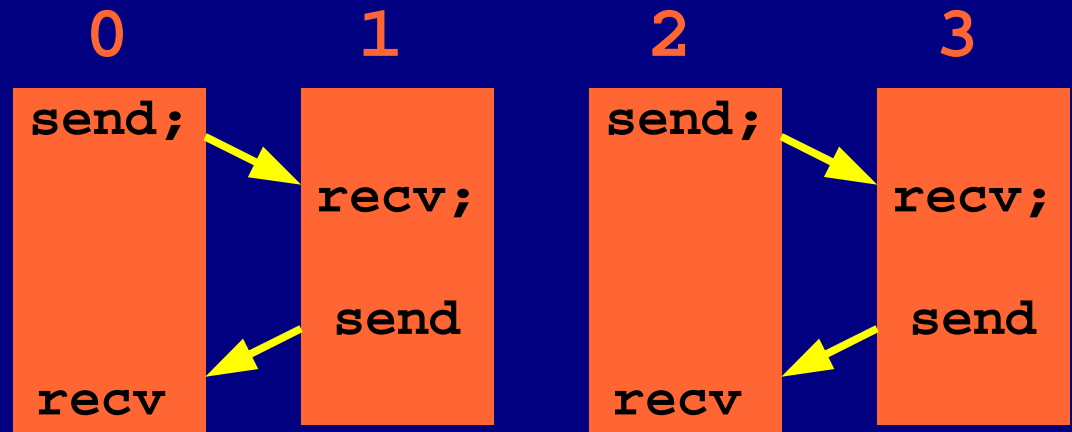


many object-programs:
one for each process

simple example

```
if even(my_proc_id)
  then .< send;
        rcv >.
  else .< rcv;
        send >.
```

`my_proc_id =`



Meta-Programming Actions

- combination of atomic parts of the specification

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- removal of interpretation overhead
 - analysis of specification

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 - case distinctions and addressing calculations for
 - process identifier
 - block of distributed data

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- removal of interpretation overhead
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- domain-specific optimization

Specification Language

(1) atomic computation: **Atom** **f**

→ **f** is a function in the host language

Specification Language

(1) atomic computation: **Atom f**

(2) sequential composition: **Seq(n, f)**

→ **n** : the number of parts in the sequence

→ **f** : mapping from index **i** to part **$f(i)$**

Specification Language

(1) atomic computation: **Atom** f

(2) sequential composition: **Seq**(n, f)

(3) parallel composition: **Par**(n, f)

→ n : the number of parallel parts

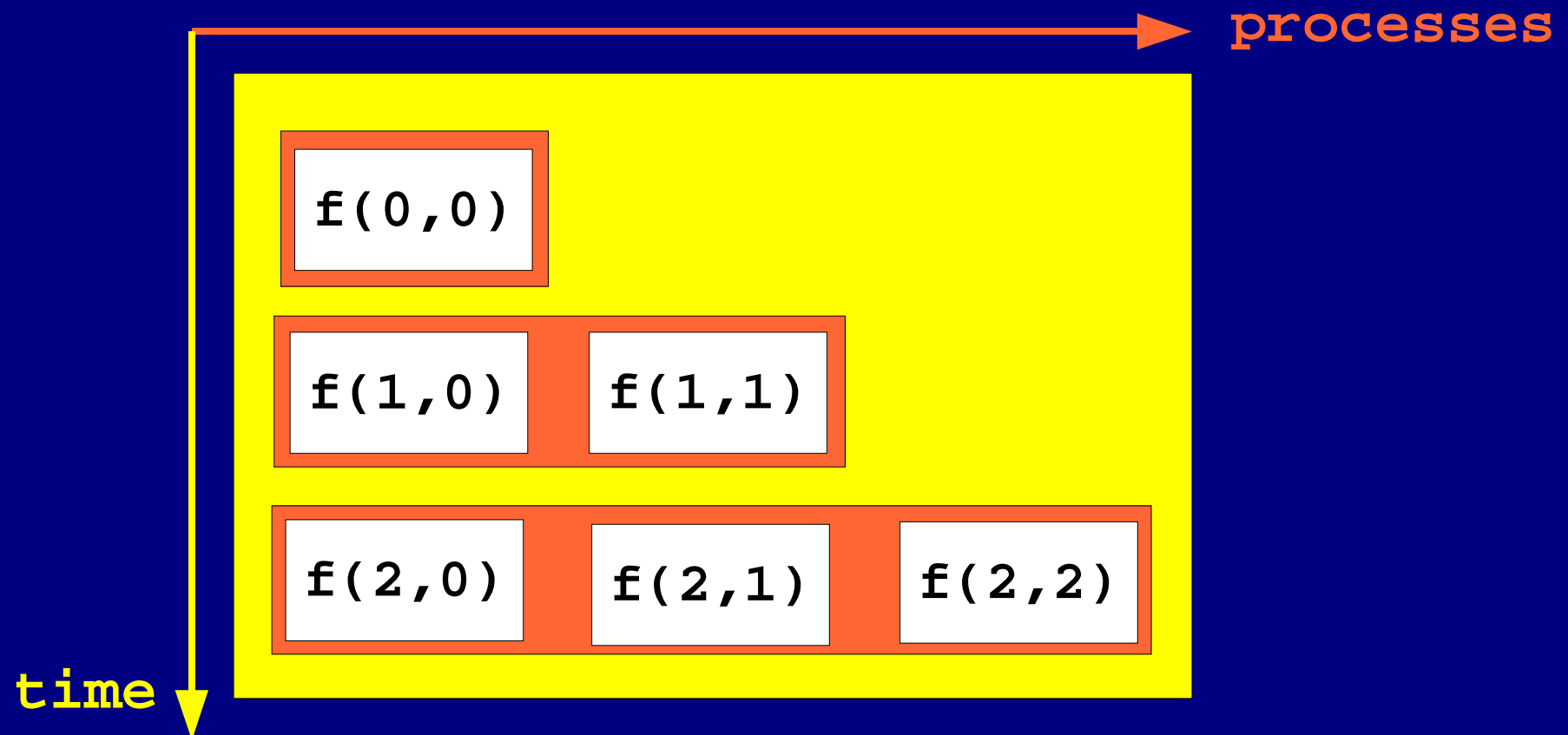
→ f : mapping from index i to part $f(i)$

Specification Example

```
let single s p = Atom (f (s,p)) in
```

```
let step s = Par (s+1, fun _ p -> single s p)
```

```
in Seq (3, fun s -> step s)
```

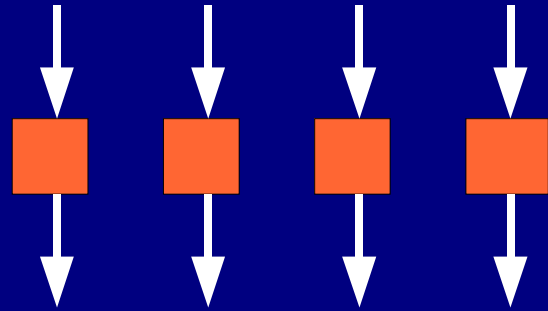


Cost Model

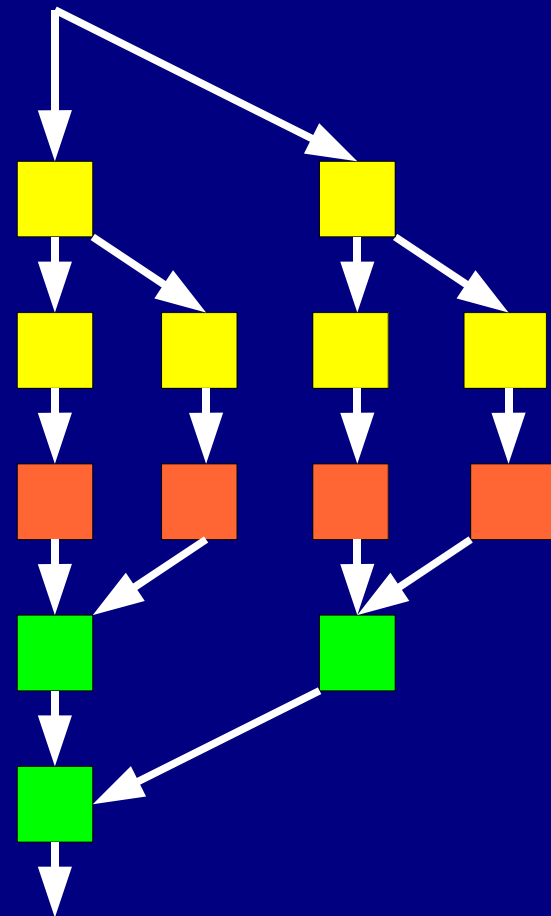
	w: work	d: depth	u: used PEs
Atom f	1	1	1
Seq(n, f)	$\sum_{0 \leq i < n} w(fi)$	$\sum_{0 \leq i < n} d(fi)$	$\max_{0 \leq i < n} u(fi)$
Par(n, f)	$\sum_{0 \leq i < n} w(fi)$	$\max_{0 \leq i < n} d(fi)$	$\sum_{0 \leq i < n} u(fi)$

Skeletons

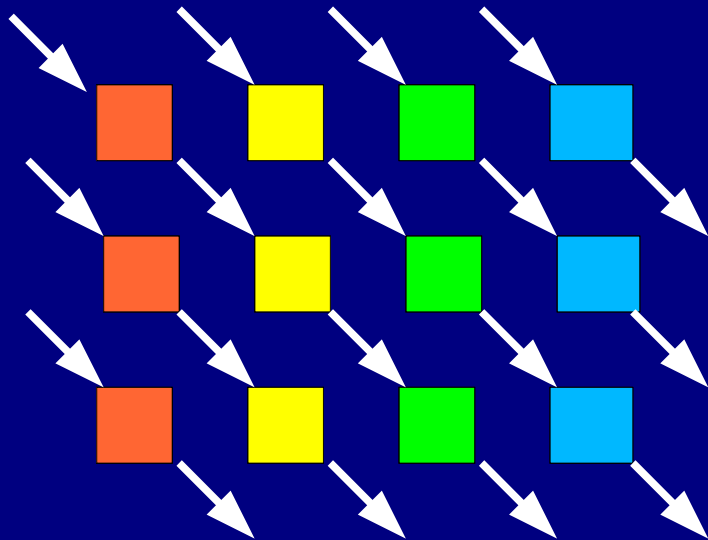
map



divide & conquer



pipeline/systolic

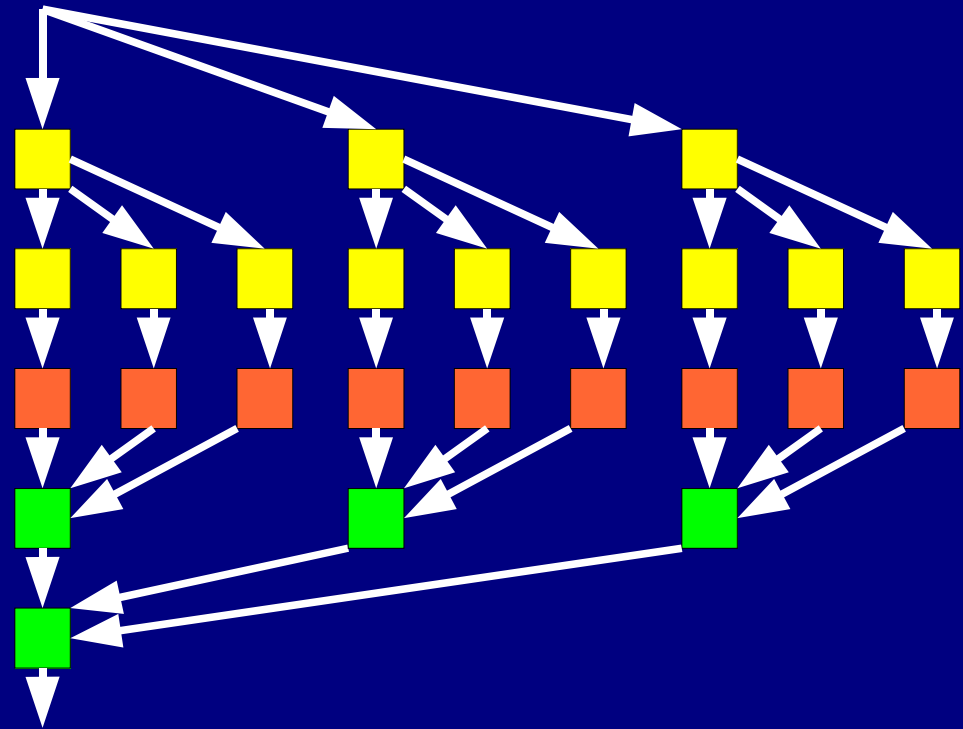


Divide & Conquer (1)

example

- recursion depth = 2
- division degree = 3

```
Par (3, ...  
  Seq (_, ...  
    Par (3, ...  
      Seq (_, ...)  
    ))  
  )  
)
```

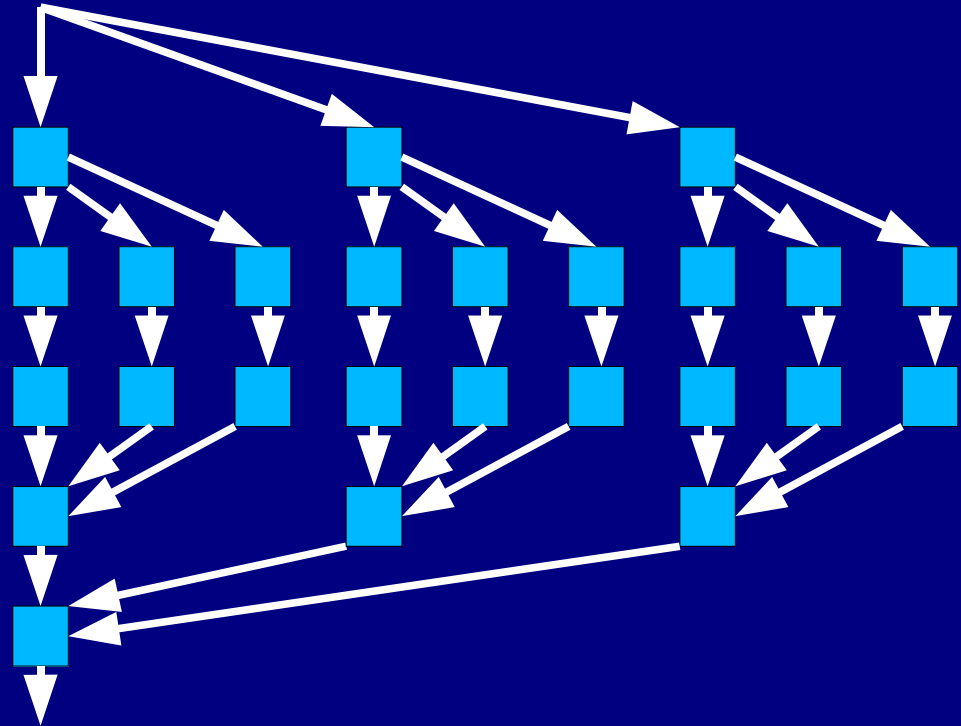


Divide & Conquer (2)

example

- recursion **depth** = 2
- division **degree** = 3

```
Par (3, ...  
  Seq (_, ...  
    Par (3, ...  
      Seq (_, ...)  
    )))
```



meta-program: recursive construction rule

dc 0 = Atom ...

dc (depth+1) = Par (degree, ...Seq(, ...dc (depth)...)...)

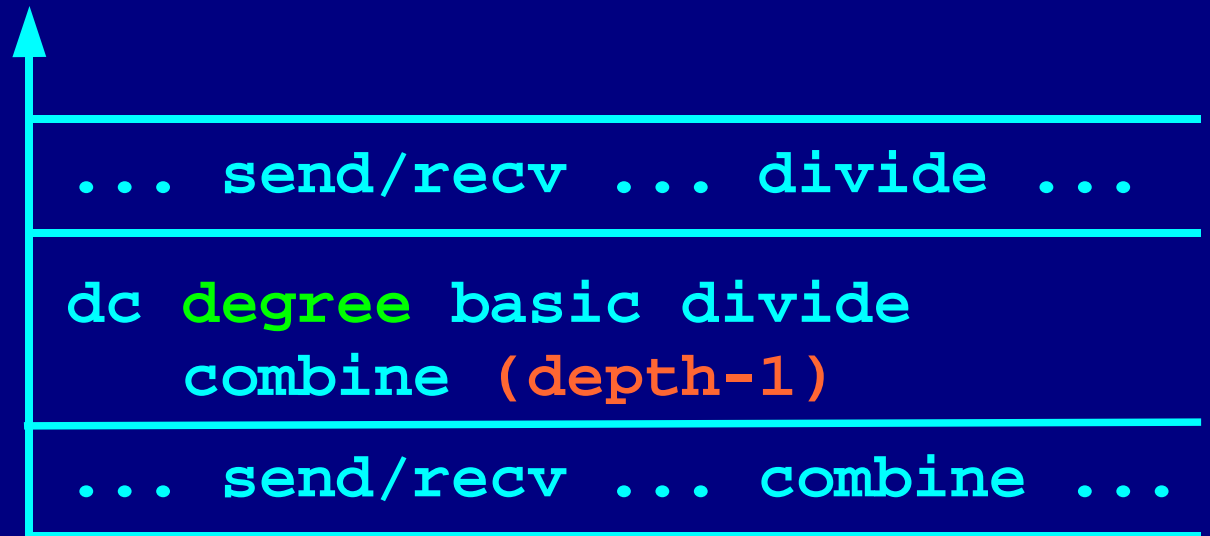
Divide & Conquer (3)

- structural parameters (influence parallelism)
 - **depth**
 - **degree**
- customizing functions
 - **basic**
 - **divide**
 - **combine**

Divide & Conquer (4)

meta-program

```
let rec dc degree basic divide combine depth =  
  if depth=0  
  then Atom (fun x -> (.< let y = .~x  
                           in basic y  
                           >.)  
  else  
  Par (degree, subtask)
```



Divide & Conquer (4a)

```
let rec dc degree basic divide combine depth =
  if depth=0
  then Atom (fun x -> (.< let y = .~x
                          in basic y
                          >.) )
  else
  Par (degree, (* subtask *)
       fun partners mypart
       -> if mypart=0
          then master partners
          else worker mypart partners.(0))
```

Divide & Conquer (4b)

```
master partners = cseq
[ Atom (fun x ->
    .< let y = .~x in
      for i=1 to degree-1 do
        send y (partners.(i)) depth
      done;
      divide 0 y
    >.);
  dc degree basic divide combine (depth-1);
  Atom (fun x ->
    .< let y = .~x in
      let tmpdata = Array.make degree [||] in
        tmpdata.(0) <- y;
        for i=1 to degree-1 do
          tmpdata.(i) <- receive (partners.(i))
                           depth
        done;
        combine tmpdata
    >.)
]
```

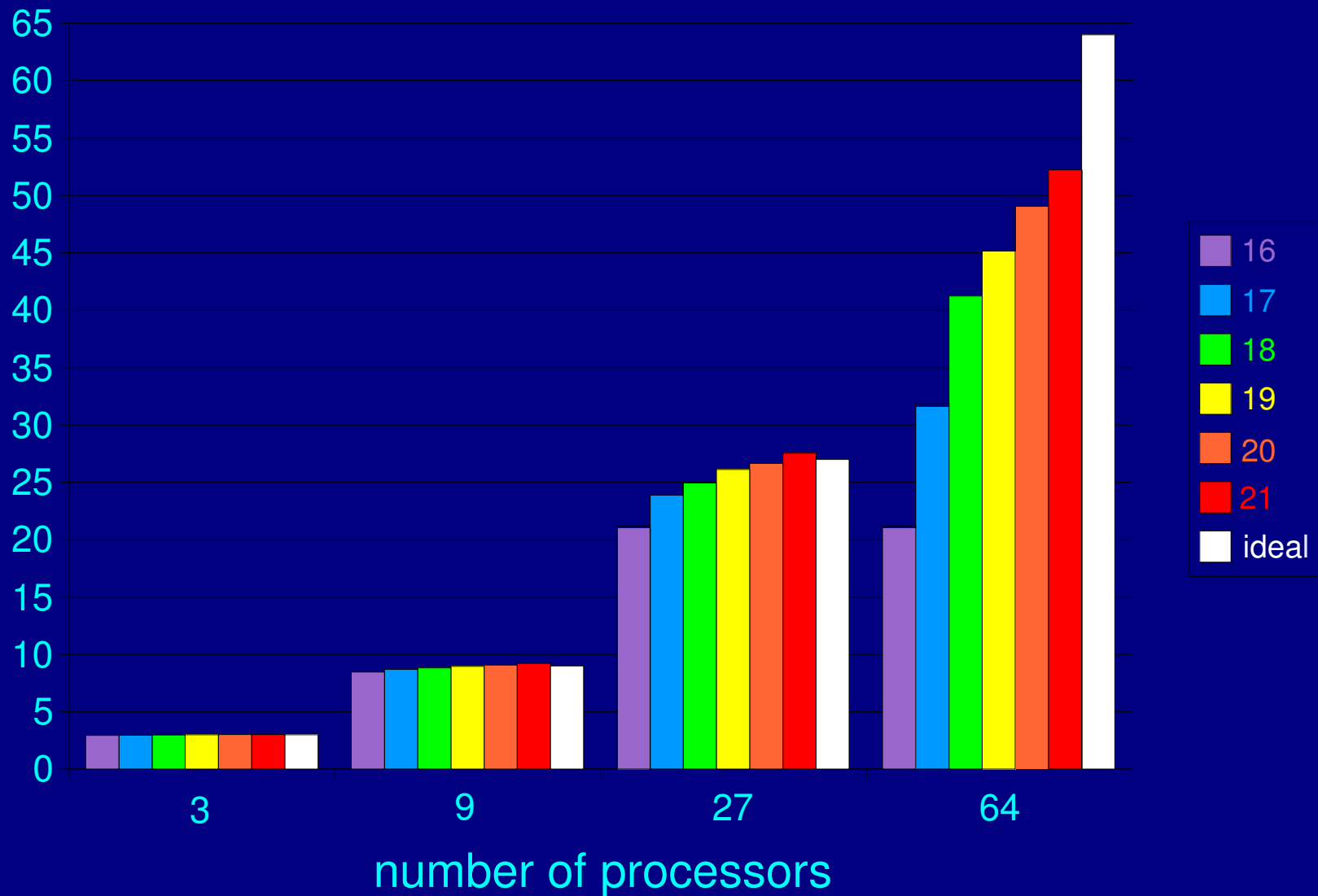
Divide & Conquer (4c)

```
worker mypart partner = cseq
[ Atom (fun x ->
    .< begin
        .~x;
        let y = receive partner depth in
            divide mypart y
        end
    >.);
dc degree basic divide combine (depth-1);
Atom (fun x ->
    .< let y = .~x in
        send y partner depth;
        y
    >.)
]
```

Experimental Results

D&C (Karatsuba Algorithm)

speedup



Experimental Results

- abstraction penalties
 - partial evaluation: **negligible** (few ms)
 - bytecode interpretation: **serious** (factor 3-8)

Conclusions

- functional meta-programming: **useful for**
 - organization of collective communications
 - definitions of parallel skeletons
- partial evaluation time is **negligible!**
- improvements of sequential parts **necessary**
 - short-term:
 - linking with compiled native code
 - call to external functions in C/Fortran
 - long-term:
 - just-in-time compilation

sources soon available via:

www.fmi.uni-passau.de/~herrmann

thank you for your attention!

questions?